

# Stochastic Linear Programming

Operations Research

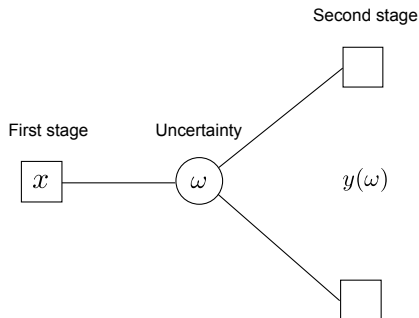
Anthony Papavasiliou

- 1 Two-Stage Stochastic Linear Programs
- 2 Scenario Trees, Lattices, and Serial Independence
- 3 Multi-Stage Stochastic Linear Programs
- 4 Applying Dynamic Programming to Stochastic Linear Programs

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# Sequence of Events

- 1 **First-stage decisions:** decisions taken before uncertainty is revealed
- 2 **Second-stage decisions:** decisions taken after uncertainty is revealed
- 3 Sequence of events:  $x \rightarrow \omega \rightarrow y(\omega)$



$$\min c^T x + \mathbb{E}[\min q(\omega)^T y(\omega)]$$

$$Ax = b$$

$$T(\omega)x + W(\omega)y(\omega) = h(\omega)$$

$$x \geq 0, y(\omega) \geq 0$$

- First-stage decisions  $x \in \mathbb{R}^{n_1}$ , second stage decisions  $y(\omega) \in \mathbb{R}^{n_2}$
- First-stage parameters:  $c \in \mathbb{R}^{n_1}$ ,  $b \in \mathbb{R}^{m_1}$ ,  $A \in \mathbb{R}^{m_1 \times n_1}$
- Second-stage data:  $q(\omega) \in \mathbb{R}^{n_2}$ ,  $h(\omega) \in \mathbb{R}^{m_2}$ ,  $T(\omega) \in \mathbb{R}^{m_2 \times n_1}$ ,  
 $W(\omega) \in \mathbb{R}^{m_2, n_2}$
- **Fixed recourse** if  $W$  does not depend on  $\omega$

# Example: Newsboy Problem

Denote

- $x$ : amount of product produced in period 1
- $y$ : amount of product sold in period 2
- $C$ : unit cost of production
- $P$ : sale price
- $D(\omega)$ : random demand

Two-stage stochastic formulation of newsboy problem:

$$\begin{aligned} \min_{x, s(\omega) \geq 0} \quad & C \cdot x - \mathbb{E}[P \cdot s(\omega)] \\ \text{s.t.} \quad & s(\omega) \leq x \\ & s(\omega) \leq D(\omega) \end{aligned}$$

Extensions: salvage value, penalty for unserved demand

What is the trade-off of large/small value of  $x$ ?

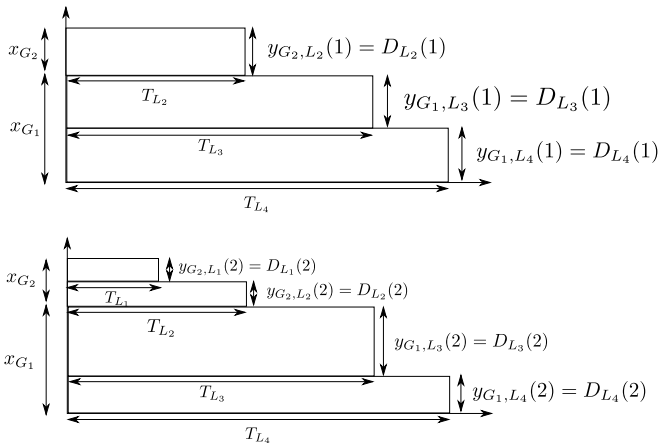
# Example: Capacity Expansion Planning

$$\begin{aligned} \min_{x, y \geq 0} & \sum_{i=1}^n (I_i \cdot x_i + \mathbb{E}[\sum_{j=1}^m C_i \cdot T_j \cdot y_{ij}(\omega)]) \\ \text{s.t.} & \sum_{i=1}^n y_{ij}(\omega) = D_j(\omega), j = 1, \dots, m \\ & \sum_{j=1}^m y_{ij}(\omega) \leq x_i, i = 1, \dots, n - 1 \end{aligned}$$

- $I_i, C_i$ : fixed/variable cost of technology  $i$
- $D_j(\omega), T_j$ : height/width of load block  $j$
- $y_{ij}(\omega)$ : capacity of  $i$  allocated to  $j$
- $x_i$ : capacity of  $i$

Note:  $T_j$  independent of  $\omega$

# Example: Capacity Expansion Planning - Graphical Illustration



Note:  $T_j$  independent of  $\omega$



# Example: Hydro-Thermal Scheduling

Denote:

- $q_t$ : hydro power
- $p_t$ : thermal power
- $C$ : marginal cost of thermal power plant
- $D_t$ : demand
- $E$ : storage limit in the dam
- $x_t$ : content of dam at the *end* of a stage
- $r_t$ : amount of rain during stage  $t$

Hydro-thermal scheduling problem:

$$\min C \cdot p_1 + \mathbb{E}[C \cdot p_2(\omega)]$$

$$p_1 + q_1 \geq D_1$$

$$x_1 \leq x_0 + r_1 - q_1$$

$$x_1 \leq E$$

$$p_2(\omega) + q_2(\omega) \geq D_2$$

$$q_2(\omega) \leq x_1 + r_2(\omega)$$

$$p_1, q_1, x_1, p_2(\omega), q_2(\omega) \geq 0$$

What is the trade-off?

- 1 Two-Stage Stochastic Linear Programs
- 2 Scenario Trees, Lattices, and Serial Independence**
- 3 Multi-Stage Stochastic Linear Programs
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A **scenario tree** is a graphical representation of a Markov process  $\{\xi_t\}_{t \in \mathbb{Z}}$ , where

- nodes correspond to histories of realizations  $\xi_{[t]} = (\xi_1, \dots, \xi_t)$
- edges correspond to transitions from  $\xi_{[t]}$  to  $\xi_{[t+1]}$

# Scenario Tree Terminology

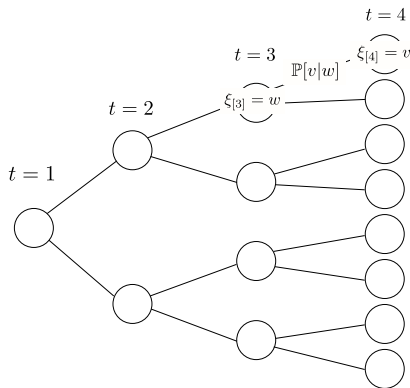
- Root corresponds to  $t = 1$
- **Ancestor** of a node  $\xi_{[t]}$ ,  $A(\xi_{[t]})$ : *unique* adjacent node which precedes  $\xi_t$ :

$$A(\xi_{[t]}) = \{\xi_{[t-1]} : (\xi_{[t-1]}, \xi_{[t]}) \in E\}$$

- **Children** or **descendants** of a node,  $C(\xi_{[t]})$ : set of nodes that are adjacent to  $\xi_{[t]}$  and occur at stage  $t + 1$ :

$$C(\xi_{[t]}) = \{\xi_{[t+1]} : (\xi_{[t]}, \xi_{[t+1]}) \in E\}$$

# Scenario Tree Graphical Illustration



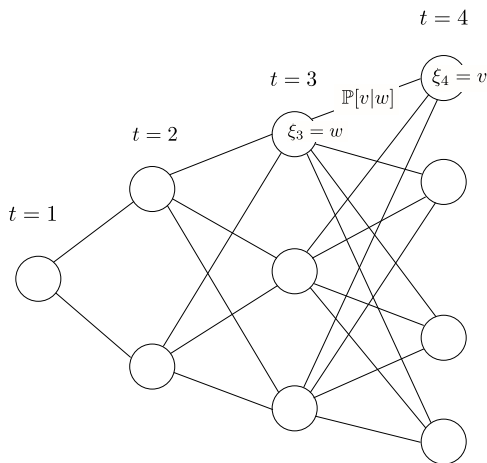
Specification of probability space requires:

- Assigning value  $\xi_{[t]}$  for every node
- Assigning value  $\mathbb{P}[\xi_{[t+1]}|\xi_{[t]}]$  for every edge

A **lattice** is a graphical representation of a Markov process  $\{\xi_t\}_{t \in \mathbb{Z}}$ , where

- nodes correspond to realizations  $\xi_t$
- edges correspond to transitions from  $\xi_t$  to  $\xi_{t+1}$

# Lattice Graphical Illustration

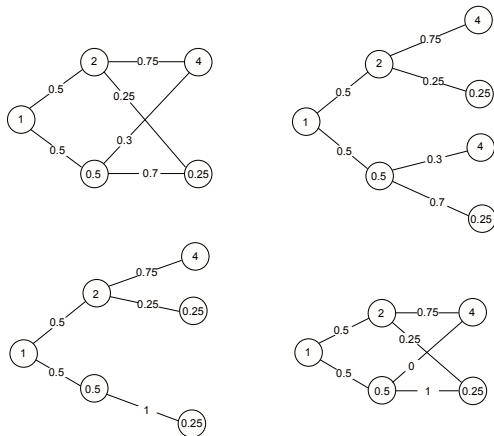


Specification of probability space requires:

- Assigning value  $\xi_t$  for every node
- Assigning value  $\mathbb{P}[\xi_{t+1}|\xi_t]$  for every edge



# Equivalence of Scenario Trees and Lattices



We can

- unfold lattices into scenario trees (top)
- fold scenario trees into lattices (bottom)

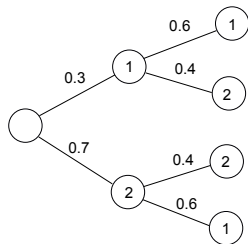
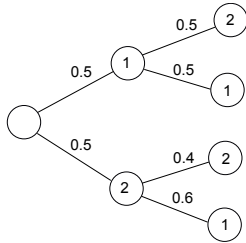
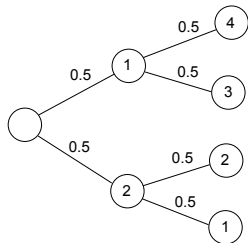
# Serial Independence

A process satisfies **serial independence** if, for every stage  $t$ ,  $\xi_t$  has a probability distribution that does not depend on the history of the process, i.e. one can define a probability measure  $p_t(i)$  at each stage  $t$ , such that

$$\mathbb{P}[\xi_t(\omega) = i | \xi_{[t-1]}(\omega)] = p_t(i), \forall \xi_{[t-1]} \in \Xi_{[t-1]}, i \in \Xi_t$$

# Checking for Serial Independence

Values on arcs indicate transition probabilities, values in nodes indicate realization of  $\xi_t$



Which scenario tree(s) obey(s) serial independence

# Populating Scenario Trees and Lattices with Data

- For scenario trees, one specifies:
  - The value of  $\xi_t$  at each node
  - The transition probability for every edge
- For lattices, one specifies:
  - The value of  $\xi_t$  at each node (a node generally does not correspond to a unique history  $\xi_{[t]}$ )
  - The transition probability for every edge
- For lattices with stage-wise independence, one specifies:
  - The value of  $\xi_t$  at each node
  - The probability of realization of each node of the lattice (well-defined)

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**Extended form** of a multistage stochastic linear program:

(*MSLP*) :

$$\min c_1^T x_1 + \mathbb{E}[c_2(\omega)^T x_2(\omega) + \dots + c_H(\omega)^T x_H(\omega)]$$

$$\text{s.t. } W_1 x_1 = h_1$$

$$T_1(\omega)x_1 + W_2(\omega)x_2(\omega) = h_2(\omega), \omega \in \Omega$$

⋮

$$T_{t-1}(\omega)x_{t-1}(\omega) + W_t(\omega)x_t(\omega) = h_t(\omega), \omega \in \Omega$$

⋮

$$T_{H-1}(\omega)x_{H-1}(\omega) + W_H(\omega)x_H(\omega) = h_H(\omega), \omega \in \Omega$$

$$x_1 \geq 0, x_t(\omega) \geq 0, t = 2, \dots, H$$

- Probability space  $(\Omega, 2^\Omega, \mathbb{P})$  with filtration  $\{\mathcal{A}\}_{t \in \{1, \dots, H\}}$
- $c_t(\omega) \in \mathbb{R}^{n_t}$ : cost coefficients
- $h_t(\omega) \in \mathbb{R}^{m_t}$ : right-hand side parameters
- $W_t(\omega) \in \mathbb{R}^{m_t \times n_t}$ : coefficients of  $x_t(\omega)$
- $T_{t-1}(\omega) \in \mathbb{R}^{m_t \times n_{t-1}}$ : coefficients of  $x_{t-1}(\omega)$
- $x_t(\omega)$ : set of state *and* action variables in period  $t$
- We implicitly enforce **non-anticipativity** by requiring that  $x_t$  and  $\xi_t$  are *adapted* to filtration  $\{\mathcal{A}\}_{t \in \{1, \dots, H\}}$

We now consider two specific instantiations of (MSLP):

- (MSLP-ST): stochastic programs on scenario trees
- (MSLP-L): stochastic programs on lattices

In these formulations, we will use the following notation:

- $\omega_t \in \mathcal{S}_t$  (interpretation: index in the support  $\Xi_t$  of random input  $\xi_t$ )
- $\omega_{[t]} \in \mathcal{S}_1 \times \dots \times \mathcal{S}_t$  (interpretation: index in  $\Xi_{[t]} = \Xi_1 \times \dots \times \Xi_t$ , which is the history of realizations, up to period  $t$ )



# Formulation on a Scenario Tree

(MSLP – ST) :

$$\min c_1^T x_1 + \mathbb{E}[c_2(\omega_{[2]})^T x_2(\omega_{[2]}) + \dots + c_H(\omega_{[H]})^T x_H(\omega_{[H]})]$$

$$\text{s.t. } W_1 x_1 = h_1$$

$$T_1(\omega_{[2]})x_1 + W_2(\omega_{[2]})x_2(\omega_{[2]}) = h_2(\omega_{[2]}), \omega_{[2]} \in S_1 \times S_2$$

⋮

$$T_{t-1}(\omega_{[t]})x_{t-1}(\omega_{[t-1]}) + W_t(\omega_{[t]})x_t(\omega_{[t]}) = h_t(\omega_{[t]}), \omega_{[t]} \in S_1 \times \dots \times S_t$$

⋮

$$T_{H-1}(\omega_{[H]})x_{H-1}(\omega_{[H-1]}) + W_H(\omega_{[H]})x_H(\omega_{[H]}) = h_H(\omega_{[H]}),$$

$$\omega_{[H]} \in S_1 \times \dots \times S_H$$

$$x_1 \geq 0, x_t(\omega_{[t]}) \geq 0, t = 2, \dots, H$$

(MSLP - L) :

$$\min c_1^T x_1 + \mathbb{E}[c_2(\omega_2)^T x_2(\omega_{[2]}) + \dots + c_H(\omega_H)^T x_H(\omega_{[H]})]$$

$$\text{s.t. } W_1 x_1 = h_1$$

$$T_1(\omega_2)x_1 + W_2(\omega_2)x_2(\omega_{[2]}) = h_2(\omega_2), \omega_{[2]} \in \mathbf{S}_1 \times \mathbf{S}_2$$

⋮

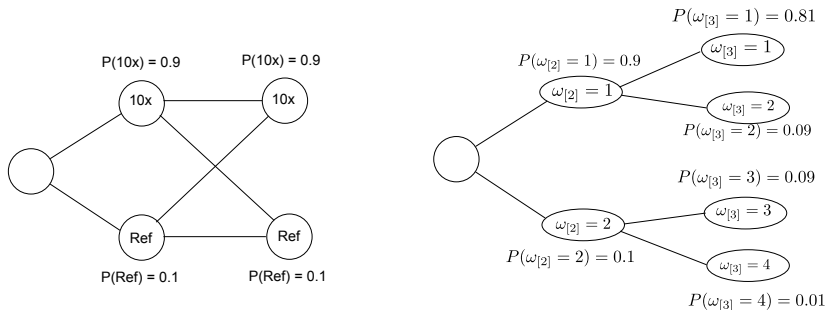
$$T_{H-1}(\omega_H)x_{H-1}(\omega_{[H]}) + W_H(\omega_H)x_H(\omega_{[H]}) = h_H(\omega_H),$$

$$\omega_{[H]} \in \mathbf{S}_1 \times \dots \times \mathbf{S}_H$$

$$x_1 \geq 0, x_t(\omega_{[t]}) \geq 0, t = 2, \dots, H$$

- Compared to (MSLP-ST),  $\xi_t$  in (MSLP-L) is indexed over  $\omega_t \in S_t$
- Problem size of (MSLP-L) doesn't really change compared to (MSLP-ST) ( $x_t$  is still indexed over  $\omega_{[t]} \in S_1 \times \dots \times S_t$ )

# Example: Capacity Expansion - Scenario Tree



**Table:** Load duration curve for reference and 10x outcome

	Duration (hours)	Level (MW)	
		Reference scenario	10x wind scenario
Base load	8760	0-7086	0-3919
Medium load	7000	7086-9004	3919-7329
Peak load	1500	9004-11169	7329-10315

# Example: Capacity Expansion - Technological Options

Technology	Fuel cost (\$/MWh)	Inv cost (\$/MWh)
Coal	25	16
Gas	80	5
Nuclear	6.5	32
Oil	160	2
DR	1000	0

# Example: Capacity Expansion - Notation and Setup

Denote:

- $v_{it\omega_{[t]}}$ : capacity of technology  $i$  constructed in period  $t$
- $x_{it\omega_{[t]}}$ : *total* amount of capacity of technology  $i$  available in period  $t$
- $y_{ijt\omega_{[t]}}$ : power allocation from technology  $i$  to load block  $j$

Sequence of events:

- 1 Capacity  $x_{i,t-1,\omega_{[t-1]}}$  available at the *end* of stage  $t - 1$  that can serve demand in  $t$
- 2 Demand  $D_{jt\omega_{[t]}}$  is observed
- 3 Construct new capacity  $v_{it\omega_{[t]}}$

# Example: Capacity Expansion - Model

Objective function:

$$\begin{aligned} \min_{x, v, y \geq 0} & \sum_{i=1}^n l_i \cdot v_{i1} \\ & + \sum_{\omega_{[2]}=1}^2 p_{\omega_{[2]}} \left( \sum_{i=1}^n l_i \cdot v_{i2\omega_{[2]}} + \sum_{i=1}^n \sum_{j=1}^m C_i \cdot T_j \cdot y_{ij2\omega_{[2]}} \right) \\ & + \sum_{\omega_{[3]}=1}^4 p_{\omega_{[3]}} \left( \sum_{i=1}^n l_i \cdot v_{i3\omega_{[3]}} + \sum_{i=1}^n \sum_{j=1}^m C_i \cdot T_j \cdot y_{ij3\omega_{[3]}} \right) \end{aligned}$$

Note: first stage involves *only* investment decision

Supply equals demand (enforced only for  $t > 1$ ):

$$\begin{aligned} \sum_{i=1}^n y_{ijt\omega_{[t]}} &= D_{jt\omega_{[t]}}, j \in \{1, \dots, m\}, t \in \{2, \dots, 3\} \\ \omega_{[2]} &\in \{1, 2\}, \omega_{[3]} \in \{1, \dots, 4\} \end{aligned}$$

# Example: Capacity Expansion - Model

Investment dynamics:

$$x_{i2\omega_{[2]}} = x_{i11} + v_{i2\omega_{[2]}}, i \in \{1, \dots, n-1\}, \omega_{[2]} \in \{1, 2\}$$

$$x_{i3\omega_{[3]}} = x_{i21} + v_{i3\omega_{[3]}}, i \in \{1, \dots, n-1\}, \omega_{[3]} \in \{1, 2\}$$

$$x_{i3\omega_{[3]}} = x_{i22} + v_{i3\omega_{[3]}}, i \in \{1, \dots, n-1\}, \omega_{[3]} \in \{3, 4\}$$

Technology capacity constraints:

$$\sum_{j=1}^m y_{ij2\omega_{[2]}} \leq x_{i11}, i \in \{1, \dots, n-1\}, \omega_{[2]} \in \{1, 2\}$$

$$\sum_{j=1}^m y_{ij3\omega_{[3]}} \leq x_{i21}, i \in \{1, \dots, n-1\}, \omega_{[3]} \in \{1, 2\}$$

$$\sum_{j=1}^m y_{ij3\omega_{[3]}} \leq x_{i22}, i \in \{1, \dots, n-1\}, \omega_{[3]} \in \{3, 4\}$$

Does this model obey block separability?



# Example: Capacity Expansion - Optimal Solution

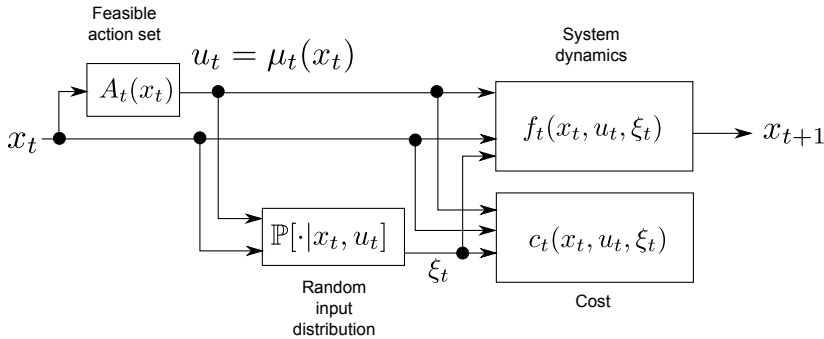
Optimal expansion plan:

- Coal, period 1: 2986 MW
- Nuclear, period 1: 7329 MW
- Oil, period 1: 854 MW
- Period 2: nothing (!)

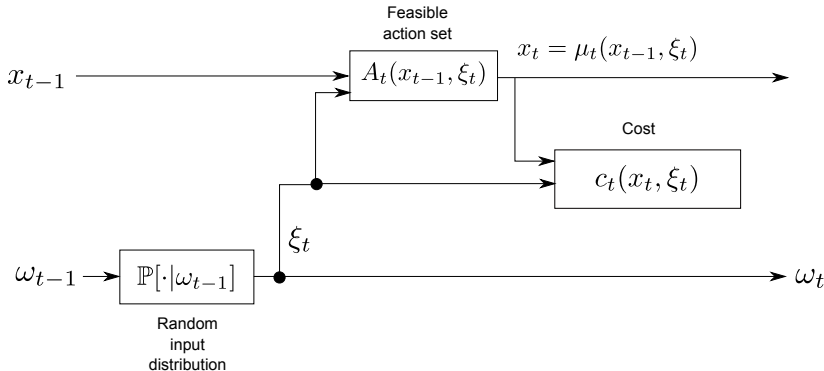
Why is it optimal to invest *only* in period 1?

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# Stochastic Control Block Diagram



# Stochastic Programming Block Diagram



- Timing of action
  - Stochastic control: *first* decide  $u_t$ , *then* observe realization of uncertainty  $\xi_t$
  - Stochastic programming: *first* observe the realization of uncertainty,  $\xi_t$ , *then* decide  $x_t$
- System state
  - Stochastic control:  $x_t$  encodes *all* information about system state
  - Stochastic programming: vector  $x_t$  *and* node of the lattice  $\omega_t$  encode all information about system state
- Feasible action set  $A_t$ 
  - Stochastic control:  $A_t$  depends only on  $x_t$
  - Stochastic programming:  $A_t$  depends on  $x_{t-1}$  *and*  $\xi_t$

Feasible action set in stage  $t$ :

$$T_{t-1}(\omega_t)x_{t-1}(\omega_{[t]}) + W_t(\omega_t)x_t(\omega_{[t]}) = h_t(\omega_t), \omega_{[t]} \in \mathcal{S}_1 \times \dots \times \mathcal{S}_t$$

**Block separability** occurs when these constraints can be written in the following form:

$$T_{t-1}^{xx}(\omega_t)x_{t-1}(\omega_{[t-1]}) + W_t^{xx}(\omega_t)x_t(\omega_{[t]}) = h_t^{xx}(\omega_t), \omega_{[t]} \in \mathcal{S}_1 \times \dots \times \mathcal{S}_t$$

$$T_{t-1}^{xu}(\omega_t)x_{t-1}(\omega_{[t-1]}) + W_t^{xu}(\omega_t)u_t(\omega_t) = h_t^{xu}(\omega_t), \omega_{[t]} \in \mathcal{S}_1 \times \dots \times \mathcal{S}_t$$

Benefit: decision variables  $u_t$  do not need to be propagated forward

**Q-function** in final period:

$$\begin{aligned} Q_H(x_{H-1}, \xi_H) = & \min_{x_H} c_H(\omega_H)^T x_H \\ & \text{s.t. } T_{H-1}(\omega_H)x_{H-1} + W_H(\omega_H)x_H = h_H(\omega_H) \\ & x_H \geq 0 \end{aligned}$$

**Value function** in final period:

$$V_H(x_{H-1}, \omega_{H-1}) = \mathbb{E}_{\xi_H}[Q_H(x_{H-1}, \xi_H)|\omega_{H-1}]$$

Proceeding recursively,  $Q$ -function in stage  $t$ :

$$\begin{aligned} Q_t(x_{t-1}, \xi_t) = & \min_{x_t} c_t(\omega_t)^T x_t + V_{t+1}(x_t, \omega_t) \\ \text{s.t. } & T_{t-1}(\omega_t)x_{t-1} + W_t(\omega_t)x_t = h_t(\omega_t) \\ & x_t \geq 0 \end{aligned}$$

Value function in stage  $t$ :

$$V_t(x_{t-1}, \omega_{t-1}) = \mathbb{E}_{\xi_t}[Q_t(x_{t-1}, \xi_t)|\omega_{t-1}]$$



Proceed backwards until:

$$\min c_1^T x_1 + V_2(x_1)$$

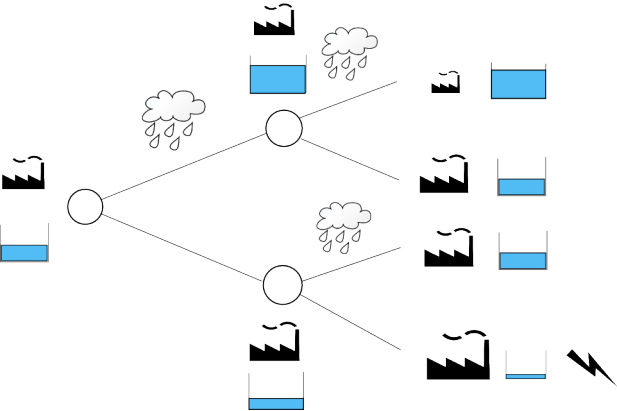
$$\text{s.t. } W_1 x_1 = h_1$$

$$x_1 \geq 0$$

# Notational Convention

- Note the different notation for node of the lattice ( $\omega_t$ ) and realization of uncertainty ( $\xi_t$ )
- The notation  $V_t(x_{t-1}, \omega_{t-1})$  emphasizes how value functions are stored by SDDP
- The notation  $Q_t(x_{t-1}, \xi_t)$  is the conventional notation used in stochastic programming, but  $Q$ -functions are not explicitly stored in SDDP
- Note the difference in the definition of the  $Q$  function
  - Stochastic control: function of state  $x$ , action  $u$
  - Stochastic programming: function of state  $x$ , random input  $\xi_t$

# Example: Hydrothermal Scheduling



# Example: Hydrothermal Scheduling

Consider the following hydro-thermal system:

- 3 periods
- Demand in each period: 1000 MW
- Marginal cost of thermal generators: 25 \$/MWh
- Max production of thermal generators: 500 MW
- Marginal cost of lost load: 1000 \$/MWh
- Rainfall: *independent* identically distributed, uniformly on  $[0, 1000]$  MW, denote density function as  $f : \mathbb{R} \rightarrow \mathbb{R}$

# Example: Hydrothermal Scheduling

Denote

- $p$ : thermal production
- $q$ : hydro production
- $l$ : unserved demand
- $x_2$ : stored hydro energy at *beginning* of period 2

$$\begin{aligned} Q_3(x_2, R_3) = & \min 1000 \cdot l + 25 \cdot p \\ & \text{s.t. } l + p + q \geq 1000 \\ & p \leq 500 \\ & q \leq x_2 + R_3 \\ & l, p, q \geq 0 \end{aligned}$$

Q function of period 3:

$$Q_3(x_2, R_3) = \begin{cases} 0, & x_2 + R_3(\omega) \geq 1000 \\ 25 \cdot (1000 - (x_2 + R_3(\omega))), & 500 \leq x_2 + R_3(\omega) < 1000 \\ 500 \cdot 25 + 1000 \cdot (500 - (x_2 + R_3(\omega))), & 0 \leq x_2 + R_3(\omega) < 500 \end{cases}$$

Value function of period 3:

$$\begin{aligned} V_3(x_2) &= \mathbb{E}_{R_3}[Q_3(x_2, R_3)] \\ &= \mathbb{P}[R_3(\omega) \geq 1000 - x_2] \cdot 0 \\ &\quad + \int_{r=500-x_2}^{1000-x_2} (25 \cdot (1000 - r - x_2))f(r)dr \\ &\quad + \int_{r=0}^{500-x_2} (500 \cdot 25 + 1000 \cdot (500 - r - x_2))f(r)dr \\ &= \begin{cases} 0, & x_2 \geq 1000 \\ 12500 - 25 \cdot x_2 + 0.0125 \cdot x_2^2, & 500 \leq x_2 < 1000 \\ 134375 - 512.5 \cdot x_2 + 0.5 \cdot x_2^2, & 0 \leq x_2 < 500 \end{cases} \end{aligned}$$

Note:

- $V_3$  is convex
- $V_3$  is *not* a piecewise linear function of  $x_2$

$Q_2$  can be computed as:

$$Q_2(x_1, R_2) = \min 1000 \cdot l + p + V_3(x_2)$$

$$\text{s.t. } l + p + q \geq 1000, p \leq 500$$

$$x_2 = x_1 - q + R_2(\omega)$$

$$l, p, q, x_2 \geq 0$$

$Q_2$  yields  $V_2$ ,  $V_2$  yields  $Q_1, \dots$