

Duality

Operations Research

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- 1 Lagrange Dual Problem
- 2 Weak and Strong Duality
- 3 Optimality Conditions
- 4 Dual Multipliers in Software

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Lagrangian Function

Standard form problem (not necessarily convex):

$$\min f_0(x)$$

$$\text{s.t. } f_i(x) \leq 0, i = 1, \dots, m$$

$$h_i(x) = 0, i = 1, \dots, p$$

$x \in \mathbb{R}^n$, domain \mathcal{D} , optimal value p^*

Lagrangian function: $L : \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^p \rightarrow \mathbb{R}$,

$\text{dom } L = \mathcal{D} \times \mathbb{R}^m \times \mathbb{R}^p$:

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

- Weighted sum of the objective and constraint functions
- λ_i is the Lagrange multiplier associated with $f_i(x) \leq 0$
- ν_i is the Lagrange multiplier associated with the equality constraint $h_i(x) = 0$

Lagrange dual function: $g : \mathbb{R}^m \times \mathbb{R}^p \rightarrow \mathbb{R}$,

$$\begin{aligned}g(\lambda, \nu) &= \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) \\ &= \inf_{x \in \mathcal{D}} (f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x))\end{aligned}$$

g is concave, can be $-\infty$ for some λ, ν

Dual Function is Concave

Consider any $(\lambda_1, \nu_1), (\lambda_2, \nu_2)$ with $\lambda_1, \lambda_2 \geq 0$ and $\alpha \in [0, 1]$

$$\begin{aligned} & g(\alpha\lambda_1 + (1 - \alpha)\lambda_2, \alpha\nu_1 + (1 - \alpha)\nu_2) \\ &= \inf_{x \in \mathcal{D}} (f_0(x) + \sum_{i=1}^m (\alpha\lambda_{1,i}f_i(x) + (1 - \alpha)\lambda_{2,i}f_i(x)) \\ & \quad + \sum_{i=1}^p (\alpha\nu_{1,i}h_i(x) + (1 - \alpha)\nu_{2,i}h_i(x))) \\ &\geq \alpha \inf_{x \in \mathcal{D}} (f_0(x) + \sum_{i=1}^m \lambda_{1,i}f_i(x) + \sum_{i=1}^p \nu_{1,i}h_i(x)) \\ & \quad + (1 - \alpha) \inf_{x \in \mathcal{D}} (f_0(x) + \sum_{i=1}^m \lambda_{2,i}f_i(x) + \sum_{i=1}^p \nu_{2,i}h_i(x)) \\ &= \alpha g(\lambda_1, \nu_1) + (1 - \alpha)g(\lambda_2, \nu_2) \end{aligned}$$

Dual Function is a Lower Bound

If $\lambda \geq 0$ then $g(\lambda, \nu) \leq p^*$

Proof: If \tilde{x} is feasible and $\lambda \geq 0$ then

$$f_0(\tilde{x}) \geq L(\tilde{x}, \lambda, \nu) \geq \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) = g(\lambda, \nu).$$

Minimizing over all feasible \tilde{x} gives $p^* \geq g(\lambda, \nu)$

Lagrange Relaxation of Stochastic Programs

Consider 2-stage stochastic program:

$$\begin{aligned} \min & f_1(x) + \mathbb{E}_\omega[f_2(y(\omega), \omega)] \\ \text{s.t.} & h_{1i}(x) \leq 0, i = 1, \dots, m_1, \\ & h_{2i}(x, y(\omega), \omega) \leq 0, i = 1, \dots, m_2 \end{aligned}$$

Introduce **non-anticipativity constraint** $x(\omega) = x$ and reformulate problem as

$$\begin{aligned} \min & f_1(x) + \mathbb{E}_\omega[f_2(y(\omega), \omega)] \\ \text{s.t.} & h_{1i}(x) \leq 0, i = 1, \dots, m_1, \\ & h_{2i}(x(\omega), y(\omega), \omega) \leq 0, i = 1, \dots, m_2 \\ (\nu(\omega)) : & x(\omega) = x \end{aligned}$$

Dual Function of Stochastic Program

$$g(\nu) = g_1(\nu) + \mathbb{E}_\omega g_2(\nu(\omega), \omega)$$

where

$$g_1(\nu) = \inf f_1(x) + \left(\sum_{\omega \in \Omega} \nu(\omega) \right)^T x$$

s.t. $h_{1i}(x) \leq 0, i = 1, \dots, m_1,$

and

$$g_2(\nu, \omega) = \inf f_2(y(\omega), \omega) - \nu x(\omega)$$

s.t. $h_{2i}(x(\omega), y(\omega), \omega) \leq 0, i = 1, \dots, m_2$

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Lagrange dual problem:

$$d^* = \max g(\lambda, \nu)$$

$$\text{s.t. } \lambda \geq 0$$

- Finds best lower bound on p^* from Lagrangian dual function
- Convex optimization problem with optimal value d^*
- λ, ν are dual feasible if $\lambda \geq 0, (\lambda, \nu) \in \text{dom } g$

Weak and Strong Duality

Weak duality: $d^* \leq p^*$

- Always holds (for convex and non-convex problems)
- Can be used for finding non-trivial bounds to difficult problems

Strong duality: $p^* = d^*$

- Does not hold in general
- Usually holds for convex problems
- Conditions that guarantee strong duality in convex problems are called constraint qualifications

Linear Programming Duality Mnemonic Table

Primal	Minimize	Maximize	Dual
Constraints	$\geq b_i$ $\leq b_i$ $= b_i$	≥ 0 ≤ 0 Free	Variables
Variables	≥ 0 ≤ 0 Free	$\leq c_j$ $\geq c_j$ $= c_j$	Constraints

Prove the mnemonic table using Lagrangian relaxation

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Complementary Slackness

If strong duality holds, x^* primal optimal, λ^*, ν^* dual optimal

$$\begin{aligned} f_0(x^*) = g(\lambda^*, \nu^*) &= \inf_x (f_0(x) + \sum_{i=1}^m \lambda_i^* f_i(x) + \sum_{i=1}^p \nu_i^* h_i(x)) \\ &\leq f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^p \nu_i^* h_i(x^*) \\ &\leq f_0(x^*) \end{aligned}$$

Therefore, the two inequalities above hold with equality and

- x^* minimizes $L(x, \lambda^*, \nu^*)$
- $\lambda_i^* f_i(x^*) = 0$ for $i = 1, \dots, m$

This is known as **complementary slackness**:

$$\lambda_i^* > 0 \Rightarrow f_i(x^*) = 0 \quad f_i(x^*) < 0 \Rightarrow \lambda_i^* = 0$$

KKT conditions for a problem with differentiable f_i, h_i :

- Primal constraints: $f_i(x) \leq 0, i = 1, \dots, m,$
 $h_i(x) = 0, i = 1, \dots, p$
- Dual constraints: $\lambda \geq 0$
- Complementary slackness: $\lambda_i f_i(x) = 0, i = 1, \dots, m$
- Gradient of the Lagrangian function with respect to x vanishes:

$$\nabla f_0(x) + \sum_{i=1}^m \lambda_i \nabla f_i(x) + \sum_{i=1}^p \nu_i \nabla h_i(x) = 0$$

From previous slide, if strong duality holds and x, λ, ν are optimal, then they must satisfy the KKT conditions

KKT Conditions of Maximization with Linear Constraints

Consider a maximization problem with linear constraints:

$$\max f(x, y)$$

$$(\lambda) : Ax + By \leq b$$

$$(\mu) : Cx + Dy = d$$

$$(\lambda_2) : x \geq 0$$

Then the KKT conditions have the following form:

$$Cx + Dy - d = 0$$

$$0 \leq \lambda \perp Ax + By - b \leq 0$$

$$0 \leq x \perp \lambda^T A + \mu^T C - \nabla_x f(x, y)^T \geq 0$$

$$\lambda^T B + \mu^T D - \nabla_y f(x, y)^T = 0$$

Example: The Diet Problem - KKT Conditions

Consider the diet problem with $b_1 = 1$ and $b_2 = 2$:

$$\min x_1 + 2x_2 + x_3$$

$$(\pi_1) : 0.5x_1 + 4x_2 + x_3 = 1$$

$$(\pi_2) : 2x_1 + x_2 + 2x_3 = 2$$

$$x_1, x_2, x_3 \geq 0$$

KKT conditions:

$$0.5x_1 + 4x_2 + x_3 = 1 \quad (1)$$

$$2x_1 + x_2 + 2x_3 = 2 \quad (2)$$

$$0 \leq x_1 \perp 0.5\pi_1 + 2\pi_2 + 1 \geq 0 \quad (3)$$

$$0 \leq x_2 \perp 4\pi_1 + \pi_2 + 2 \geq 0 \quad (4)$$

$$0 \leq x_3 \perp \pi_1 + 2\pi_2 + 1 \geq 0 \quad (5)$$

Example: The Diet Problem - Equivalent KKT Systems

Note: Since $h(x) = 0 \Leftrightarrow -h(x) = 0$, we can get a different KKT system with a different solution π^* :

$$0.5x_1 + 4x_2 + x_3 = 1 \quad (6)$$

$$2x_1 + x_2 + 2x_3 = 2 \quad (7)$$

$$0 \leq x_1 \perp -0.5\pi_1 - 2\pi_2 + 1 \geq 0 \quad (8)$$

$$0 \leq x_2 \perp -4\pi_1 - \pi_2 + 2 \geq 0 \quad (9)$$

$$0 \leq x_3 \perp -\pi_1 - 2\pi_2 + 1 \geq 0 \quad (10)$$

Claim:

- Primal optimal solution: $(x^*)^T = (0, 0.1429, 0.4286, 0)$
- Dual optimal solution: $(\pi^*)^T = (0.4286, 0.2857)$

Proof: verify that x^* and π^* satisfy KKT conditions

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Non-Uniqueness of KKT Conditions

- 1 The KKT conditions of a problem depend on how we define the Lagrangian function
- 2 The sign of dual multipliers depends on the KKT conditions (therefore, how we define the Lagrangian function)
- 3 The sensitivity interpretation of dual multipliers depends on the KKT conditions (therefore, how we define the Lagrangian function)
- 4 Different software interprets user syntax differently!



Dual Multipliers in AMPL

In order to be able to anticipate the sign of multipliers that AMPL will assign to constraints, note that:

- A constraint of the form $f_1(x) \leq, =, \geq f_2(x)$ is equivalently expressed as $f_1(x) - f_2(x) \leq, =, \geq 0$,
- the constraints are relaxed by subtracting their product with their corresponding multiplier from the Lagrangian function,
- the sign of the dual multiplier is such that the Lagrangian function provides a bound to the optimization problem,
- the primal-dual optimal pair is such that the KKT conditions corresponding to this Lagrangian function are satisfied.
- In this way, the dual multipliers reported by AMPL can always be interpreted as sensitivities.

Example

$$\{\min x + 2y \text{ s.t. } 0 \leq x, (\lambda_1), x \leq 2, (\lambda_2), y = 1, (\mu)\}$$

Objective function $f(x, y) = x + 2y$, inequality constraints

$f_1(x, y) = -x \leq 0$ (i.e., $a \leq$ constraint), $f_2(x, y) = x - 2$,

$h(x, y) = y - 1$

- AMPL Lagrangian:

$$L(x, y) = (x + 2y) - \lambda_1(-x) - \lambda_2(x - 2) - \mu(y - 1)$$

KKT conditions:

- Primal feasibility: $g_1(x, y) \leq 0, g_2(x, y) \leq 0, h(x, y) = 0$
- Dual feasibility: $\lambda_1 \leq 0, \lambda_2 \leq 0$
- Complementary: $\lambda_1 \perp g_1(x, y), \lambda_2 \perp g_2(x, y)$
- Stationarity:

$$\nabla f(x, y) - \lambda_1 \nabla g_1(x, y) - \lambda_2 \nabla g_2(x, y) - \mu \nabla h(x, y) = 0$$

Solution: $x = 0, y = 1, \lambda_1 = -1, \lambda_2 = 0, \mu = 2$